FINITE-VALUED STREAMING STRING TRANSDUCERS

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Finite-valued regular relations form a class of binary relations that enjoys good algorithmic properties



Theorem: We can decide in polynomial space whether a given SST defines a finite-valued regular relation

Theorem: Every finite-valued regular relation can be decomposed into a finite union of regular functions



Transducers are abstract machines that recognise relations

$\mathbf{R} \subset \Sigma^* \times \Gamma^*$

Rational relations recognised by finite state transducers

Regular relations recognised by streaming string transducers

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Regular relations

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Finite state transducer (FST)	
	$\mathbf{a} \epsilon \mathbf{b} \mathbf{b} \mathbf{a} \epsilon$
Input: a b b a a a b b b a b a a a	$\tau \cdot \mathcal{A}^{a \epsilon} \mathcal{A}^{a \epsilon} \mathcal{A}$
Output: bb b bb	
T_1 produces one of the b-blocks of its in	$b \epsilon$ ' $b \epsilon$ put

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	$\mathbf{a} \epsilon \mathbf{b} \mathbf{b} \mathbf{a} \epsilon$
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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbbabaaa	$T_{a} \rightarrow B_{a} B_{a}$
R_1 : a R_2 : b	
T_2 sorts its input	b R ₂ ≔ R ₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbbabaaa	$T_{a} \rightarrow B_{c} B_{c}$
R_1 : a R_2 : b b	
T_2 sorts its input	b R₂ ≔ R₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbbabaaa	$T_{a} \rightarrow \bigoplus B_{1}B_{2}$
R_1 : a R_2 : b b	
T_2 sorts its input	b R₂ ≔ R₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaa <mark>abbbabaaa</mark>	$T_{2} \rightarrow \bigoplus B_{1}B_{2}$
R_1 : a a a R_2 : b b	
T_2 sorts its input	b H₂ ≔ H₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbbabaaa	$T_{a} \rightarrow A \to B_{1}B_{2}$
R_1 : a a a a R_2 : b b	
T_2 sorts its input	b R₂ ≔ R₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbbaaaa	$T_{2} \rightarrow \bigoplus B_{1}B_{2}$
R_1 : a a a R_2 : b b b	
T_2 sorts its input	b H₂ ≔ H₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
	$T_2: \rightarrow \rightarrow \mathbb{R}_1\mathbb{R}_2$
$R_1: [a]a]a]a R_2: [b]b]b]b$	1) b R₂ ≔ R₂ b
T_2 sorts its input	

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$a \mid R_1 \coloneqq R_1 a$
$T_2: \rightarrow \bigcirc \rightarrow R_1R_2$
$\mathfrak{b} \mid R_2 \coloneqq R_2 \mathfrak{b}$

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Streaming string transducer (SST)	
	a R₁ ≔ R₁a
	$T_2: \rightarrow \mathbb{A} \to \mathbb{R}_1\mathbb{R}_2$
	b R ₂ ≔ R ₂ ъ
I ₂ sorts its input	

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
	$T_2: \rightarrow \bigcirc \rightarrow \mathbb{R}_1\mathbb{R}_2$
T_2 sorts its input	b R₂ ≔ R₂ b

Rational relations recognised by finite state transducers

Regular relations red

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
	$T_2: \rightarrow \bigcirc \rightarrow R_1R_2$
T_2 sorts its input	$\mathtt{b} \mid R_2 \coloneqq R_2 \ \mathtt{b}$

Rational relations red Regular relations red

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
R ₁ : a a a a a a R ₂ : b b b b b b	$T_2: \rightarrow \bigcirc \rightarrow \mathbb{R}_1\mathbb{R}_2$
T_2 sorts its input	b R₂ ≔ R₂ b

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Streaming string transducer (SST)	
	$a \mid R_1 \coloneqq R_1 a$
Input: abbaaabbaaaa	$T_{2} \rightarrow \bigoplus B_{1}B_{2}$
R ₁ : aaaaaaa R ₂ : bbbbbb	
T_2 sorts its input	b H₂ ≔ H₂ b

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Equivalence of rational relations is undecidable

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Equivalence of rational relations is undecidable

Equivalence of regular functions is decidable

 \vdash each input is mapped to at most 1 output



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recognised by streaming string transducers

functions

Equivalence of rational relations is undecidable

Equivalence of **regular functions** is **decidable**

 \mapsto each input is mapped to at most 1 output

Equivalence of **finite-valued regular relations** is **decidable** $\downarrow \exists k \in \mathbb{N} \text{ s.t. each input is mapped to at most k outputs}$



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- \Rightarrow The equivalence problem for finite-valued SST is in ELEMENTARY
- ⇒ Finite-valued 2-way FST are as expressive as finite-valued SST



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Raised as open problems by [2011. Alur, Deshmukh]

Known to hold for FST [1989. Weber], [1993. Weber];

...and for SST with a single register [2017. Gallot et al.]



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DECIDING FINITE VALUEDNESS

Theorem: We can decide in polynomial space whether a given SST defines a finite-valued regular relation

Forbidden pattern: (inspired by [2008. De Souza])



- The relation recognised by the pattern is not 1-valued
- · The substitutions produced on the loops have idempotent structure

Theorem: Every finite-valued regular relation can be decomposed into a finite union of regular functions

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

 $\mathcal{T} \quad \textbf{k-valued} \qquad (at most \ \textbf{k} \ outputs \ per \ input)$ $\mathcal{V} \quad \textbf{k-ambiguous} \qquad (at most \ \textbf{k} \ runs \ per \ input)$ $\mathcal{V} \quad \mathcal{T}'_1, \dots, \mathcal{T}'_k \quad unambiguous \qquad (at most \ \textbf{1} \ run \ per \ input)$ $\mathcal{T}_1, \dots, \mathcal{T}_k \quad functional \qquad (at most \ \textbf{1} \ output \ per \ input)$

Theorem: Every finite-valued regular relation can be decomposed into a finite union of regular functions

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

 \mathcal{T} k-valued(at most k outputs per input) \checkmark \mathcal{T}' only keeps the runs of \mathcal{T} that are far from each other \mathcal{T}' k-ambiguous(at most k runs per input) \checkmark \checkmark $\mathcal{T}'_1, ..., \mathcal{T}'_k$ unambiguous(at most 1 run per input) \checkmark $\mathcal{T}_1, ..., \mathcal{T}_k$ functional(at most 1 output per input)

Theorem: Every finite-valued regular relation can be decomposed into a finite union of regular functions

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

 $\mathcal{T} \quad \mathbf{k}\text{-valued} \qquad (at most \ \mathbf{k} \text{ outputs per input})$ $\mathcal{V} \quad \mathcal{T}' \quad \mathbf{k}\text{-ambiguous} \qquad (at most \ \mathbf{k} \text{ runs per input})$ $\mathcal{V} \quad \mathcal{T}'_i \text{ copies the } i\text{th lexicographically smallest run of } \mathcal{T}'$ $\mathcal{T}'_1, \dots, \mathcal{T}'_k \quad \text{unambiguous} \qquad (at most \ \mathbf{1} \text{ run per input})$ $\mathcal{V} \quad \mathcal{T}_1, \dots, \mathcal{T}_k \quad \text{functional} \qquad (at most \ \mathbf{1} \text{ output per input})$

Theorem: Every finite-valued regular relation can be decomposed into a finite union of regular functions

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

 \mathcal{T} k-valued(at most k outputs per input) \mathcal{V} \mathcal{T}' k-ambiguous(at most k runs per input) \mathcal{V} \mathcal{T}'_1 ..., \mathcal{T}'_k unambiguous(at most 1 run per input) \mathcal{V} $\mathcal{T}_i = \mathcal{T}'_i$ \mathcal{T}_1 ..., \mathcal{T}_k functional(at most 1 output per input)

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set to begin in the academic year 2025-2026

Contact me for more information: ismael.jecker@gmail.com